

# Anti-chains of mappings from $\omega^\omega$ on some BQO

J.Duparc

Equipe de Logique Mathématique,  
CNRS URA 753 et Université Paris VII

*U.F.R. de Mathématiques, 2 place Jussieu, 75251 Paris Cedex 05, France.*

duparc@logique.jussieu.fr

**Issue:** Let  $(P, \leq_P)$  be some BQO. Louveau and Saint-Raymond [1] showed that the following structure  $(\mathcal{F}, \leq_{\mathcal{F}})$  is also a BQO:

$$\mathcal{F} = \{ \phi : \omega^\omega \mapsto P : \phi \text{ is Borel with countable image} \}$$

with the usual topology on  $\omega^\omega$  and the discrete topology on the BQO  $P$ ; and  $\phi \leq_{\mathcal{F}} \psi$  iff there exists some continuous function  $h : \omega^\omega \mapsto \omega^\omega$  such that for all  $x \in \omega^\omega$   $\phi(x) \leq_P \psi \circ h(x)$ .

The following proposition answers the question of the relation between cardinalities of anti-chains of  $P$  and anti-chains of  $\mathcal{F}$ ?

## Proposition 1

1. *Every anti-chain in  $P$  has cardinality 1  $\implies$  every anti-chain in  $\mathcal{F}$  has cardinality 1*
2. *There exists an anti-chain in  $P$  of cardinality 2, but no element of  $P$  is incomparable with two different elements  $\implies$  every anti-chain in  $\mathcal{F}$  has cardinality at most 2*
3. *There exists an element in  $P$  which is incomparable with two different elements  $\implies$  there exists anti-chains of any cardinality in  $\mathcal{F}$ .*

*Proof of proposition 1:*

1. Obvious since  $\mathbb{I}$  doesn't have a w.s. in  $\mathbb{G}_{\leq}(\phi, \psi)$  implies by determinacy that  $\mathbb{I}$  has a w.s. in  $\mathbb{G}_{\leq}(\psi, \phi)$ ; hence  $\leq_{\mathcal{F}}$  is total.
2. For any  $q \in P$  define  $\check{q}$  by

$$\begin{aligned} \check{q} &= q \text{ if } q \text{ is comparable with every other element from } P \\ &= \text{the unique element in } P \text{ that is incomparable with } q, \text{ otherwise.} \end{aligned}$$

And given  $\phi \in \mathcal{F}$  we define  $\check{\phi}$  by  $\check{\phi}(x) = \widehat{\check{\phi}(x)}$ .

We say  $\phi$  is self dual whenever  $\check{\phi} \equiv_{\mathcal{F}} \phi$  and  $\phi$  is non self dual whenever  $\check{\phi} \not\equiv_{\mathcal{F}} \phi$ . We remark

**Claim 2**  $\psi \not\leq_{\mathcal{F}} \phi \implies \check{\phi} \leq_{\mathcal{F}} \psi$

*Proof of claim 2:* Since  $\mathbb{I}$  doesn't have a w.s. in  $\mathbb{G}_{\leq}(\psi, \phi)$ , by determinacy  $I$  has a w.s. in the same game. Winning conditions for  $I$  are to produce a real  $x$  that satisfies  $\psi x >_{\mathcal{P}} \phi(y)$  or  $\psi(x) \perp \phi(y)$  against  $\mathbb{I}$ 's run  $y$ . Hence  $\mathbb{I}$  has a w.s. in  $\mathbb{G}_{\leq}(\check{\phi}, \psi)$ .  $\square$

**Claim 3**  $\phi$  self dual  $\implies \phi$  is comparable with every other mapping in  $\mathcal{F}$

*Proof of claim 3:* Immediate from claim 2.  $\square$

**Claim 4**

$$(\phi \text{ non self dual} \wedge \psi \not\leq_{\mathcal{F}} \phi) \implies (\psi \equiv_{\mathcal{F}} \check{\phi} \vee \psi \geq_{\mathcal{F}} \phi)$$

*Proof of claim 4:* Assume  $\check{\phi} \not\leq_{\mathcal{F}} \psi$ , hence  $I$  has a w.s. in  $\mathbb{G}_{\leq}(\psi, \check{\phi})$  which is a w.s. for  $\mathbb{I}$  in  $\mathbb{G}_{\leq}(\phi, \psi)$ .  $\square$

This leads to:

**Proposition 5**

$$\psi \perp \phi \implies \psi \equiv_{\mathcal{F}} \check{\phi}$$

showing that antichains in  $\mathcal{F}$  have length at most 2.

3. Given  $\phi$  and  $\psi$  we define  $\phi \rightarrow \psi$  as the mapping  $\varphi$  such that for all  $x \in (\omega \setminus \{0\})^\omega$   $\varphi(x) = \phi(x')$  - where  $x'(n) = x(n) - 1$  ( $\forall n \in \omega$ ) - and  $\varphi(x) = \psi(y)$  for all  $x = u \wedge \langle 0 \rangle \wedge y \in \omega^\omega$  where  $u(i) \neq 0$  ( $\forall i < lh(u)$ ).

- (a) Assume there exists three different elements  $a, 0, 1 \in P$  such that  $a \perp 0, a \perp 1$  and  $0 <_{\mathcal{P}} 1$ .

Define:

$$\begin{aligned} \psi_0 &= \bar{a} \rightarrow \bar{0} \rightarrow \bar{a} \rightarrow \bar{0}, \\ \psi_1 &= \bar{a} \rightarrow \bar{1}. \end{aligned}$$

And for each  $u \in 2^{<\omega} \setminus \langle \rangle$  define  $\phi_u$  by

$$\phi_u = \psi_{u(0)} \rightarrow \psi_{u(1)} \rightarrow \dots \rightarrow \psi_{u(lh(u)-1)}$$

**Proposition 6** Let  $u, v \in 2^{<\omega}$  with  $lh(u) = lh(v)$ ,

$$u \neq v \implies \phi_u \perp \phi_v$$

*Proof of proposition 6:* The proof relies on the following

**Claim 7** Given  $u$  and  $v$  same length,

$I$  has winning strategies in both  $\mathbb{G}_{\leq}(\bar{0} \rightarrow \phi_u, \phi_v)$  and  $\mathbb{G}_{\leq}(\bar{1} \rightarrow \phi_u, \phi_v)$ .

*Proof of claim 7:* The proof is by induction on  $lh(u)$

i.  $lh(u) = 1$ . We need to check 8 different games. In each of them the winning strategy for  $I$  is straightforward:

A.  $\mathbb{G}_{\leq}(\bar{0} \rightarrow \bar{a} \rightarrow \bar{0} \rightarrow \bar{a} \rightarrow \bar{0}, \bar{a} \rightarrow \bar{0} \rightarrow \bar{a} \rightarrow \bar{0})$

B.  $\mathbb{G}_{\leq}(\bar{1} \rightarrow \bar{a} \rightarrow \bar{0} \rightarrow \bar{a} \rightarrow \bar{0}, \bar{a} \rightarrow \bar{0} \rightarrow \bar{a} \rightarrow \bar{0})$

C.  $\mathbb{G}_{\leq}(\bar{0} \rightarrow \bar{a} \rightarrow \bar{1}, \bar{a} \rightarrow \bar{1})$

D.  $\mathbb{G}_{\leq}(\bar{1} \rightarrow \bar{a} \rightarrow \bar{1}, \bar{a} \rightarrow \bar{1})$

E.  $\mathbb{G}_{\leq}(\bar{0} \rightarrow \bar{a} \rightarrow \bar{0} \rightarrow \bar{a} \rightarrow \bar{0}, \bar{a} \rightarrow \bar{1})$

F.  $\mathbb{G}_{\leq}(\bar{1} \rightarrow \bar{a} \rightarrow \bar{0} \rightarrow \bar{a} \rightarrow \bar{0}, \bar{a} \rightarrow \bar{1})$

G.  $\mathbb{G}_{\leq}(\bar{0} \rightarrow \bar{a} \rightarrow \bar{1}, \bar{a} \rightarrow \bar{0} \rightarrow \bar{a} \rightarrow \bar{0})$

H.  $\mathbb{G}_{\leq}(\bar{1} \rightarrow \bar{a} \rightarrow \bar{1}, \bar{a} \rightarrow \bar{0} \rightarrow \bar{a} \rightarrow \bar{0})$

ii.  $lh(u) > 1$ .  $u = \langle i_u \rangle \hat{\ } u'$ ,  $v = \langle i_v \rangle \hat{\ } v'$ .  $\phi_u$  and  $\phi_v$  end up with  $\rightarrow \bar{0}$  or an  $\rightarrow \bar{1}$ . By induction hypothesis  $I$  has a w.s. in the following games:  $\mathbb{G}_{\leq}(\bar{0} \rightarrow \phi_{u'}, \phi_{v'})$ ,  $\mathbb{G}_{\leq}(\bar{1} \rightarrow \phi_{u'}, \phi_{v'})$ ,  $\mathbb{G}_{\leq}(\bar{0} \rightarrow \phi_{\langle i_u \rangle}, \phi_{\langle i_v \rangle})$ ,  $\mathbb{G}_{\leq}(\bar{1} \rightarrow \phi_{\langle i_u \rangle}, \phi_{\langle i_v \rangle})$ . From which we get winning strategies for  $I$  in  $\mathbb{G}_{\leq}(\bar{0} \rightarrow \phi_u, \phi_v)$ ,  $\mathbb{G}_{\leq}(\bar{1} \rightarrow \phi_u, \phi_v)$ .

7  $\square$

6  $\square$

(b) Assume there exists three different elements  $a, b, c \in P$  such that  $a \perp b \perp c \perp a$ .

Given  $u \in 2^{<\omega}$  define the sequence  $v_u \in \{a, b, c\}^{lh(u)}$  by

$$\begin{aligned} v_u(0) &= a \quad \text{iff} \quad u(0) = 0 \\ &= b \quad \text{iff} \quad u(0) = 1 \end{aligned}$$

$$\begin{aligned} v_u(i+1) &= a \quad \text{iff} \quad (u(0) = 0 \wedge v_u(i) = c) \vee (u(0) = 1 \wedge v_u(i) = b) \\ &= b \quad \text{iff} \quad (u(0) = 0 \wedge v_u(i) = a) \vee (u(0) = 1 \wedge v_u(i) = c) \\ &= c \quad \text{iff} \quad (u(0) = 0 \wedge v_u(i) = b) \vee (u(0) = 1 \wedge v_u(i) = a) \end{aligned}$$

And finally define the mapping  $\phi_u$  by

$$\phi_u = v_u(0) \rightarrow v_u(1) \rightarrow \dots \rightarrow v_u(lh(u) - 1)$$

**Claim 8** Let  $u, v \in 2^{<\omega}$  with  $lh(u) = lh(v)$ ,

$$u \neq v \implies \phi_u \perp \phi_v$$

*Proof of claim 8:* Straightforward.

8  $\square$

1  $\square$

## References

- [1] **A. Louveau and J. Saint-Raymond** [1990] *On the quasi-ordering of Borel linear orders under embeddability* J. Symbolic Logic **55** (1990), no. 2, 537–560;